



Approaching Ideal NoC Latency with Pre-Configured Routes

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Introduction

- Problem: Latency NoCs impose.
- Motivation: Latency introduced to every communication pair.
- Past work: Achieves 1 cycle/hop at 500 MHz.
- We extend speculation to routing decisions.
- Goal: Approach buffered wire latency.
 - Fraction of cycle/hop.



Our Approach

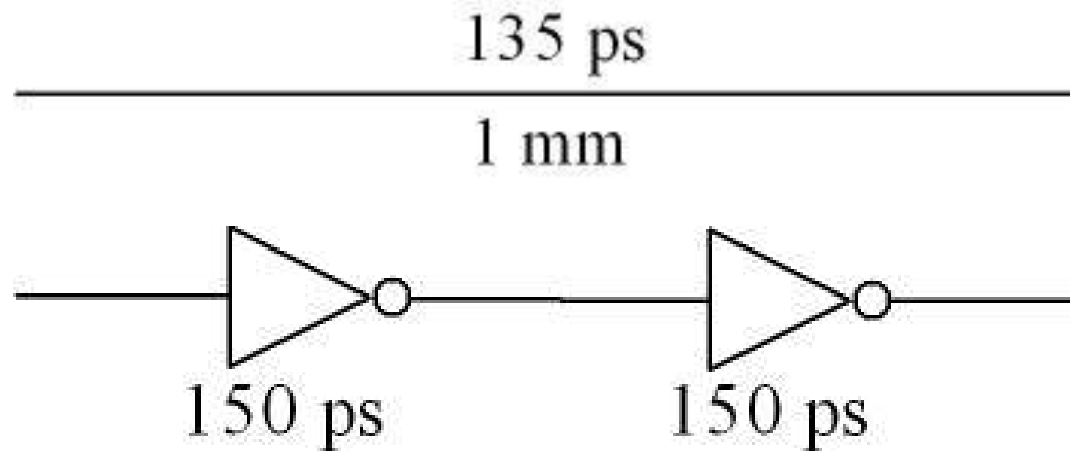
130 nm
library

135 ps

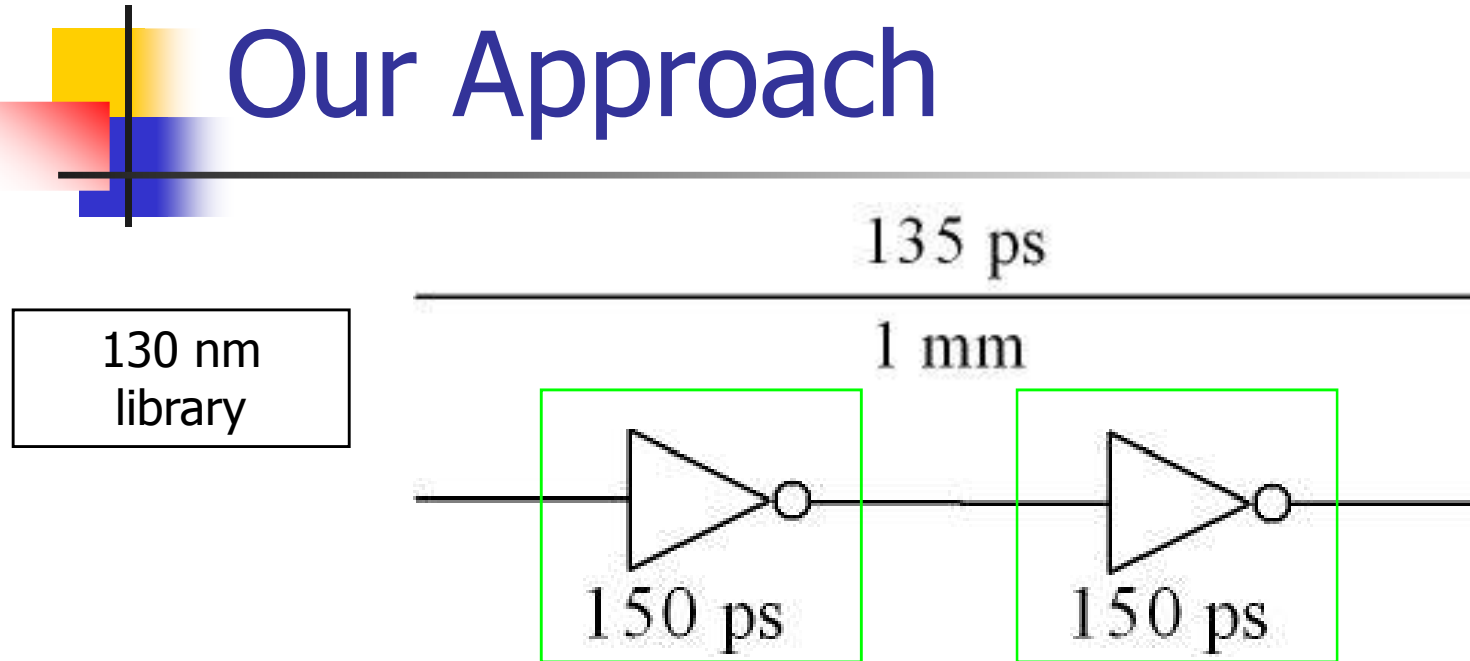
1 mm

Our Approach

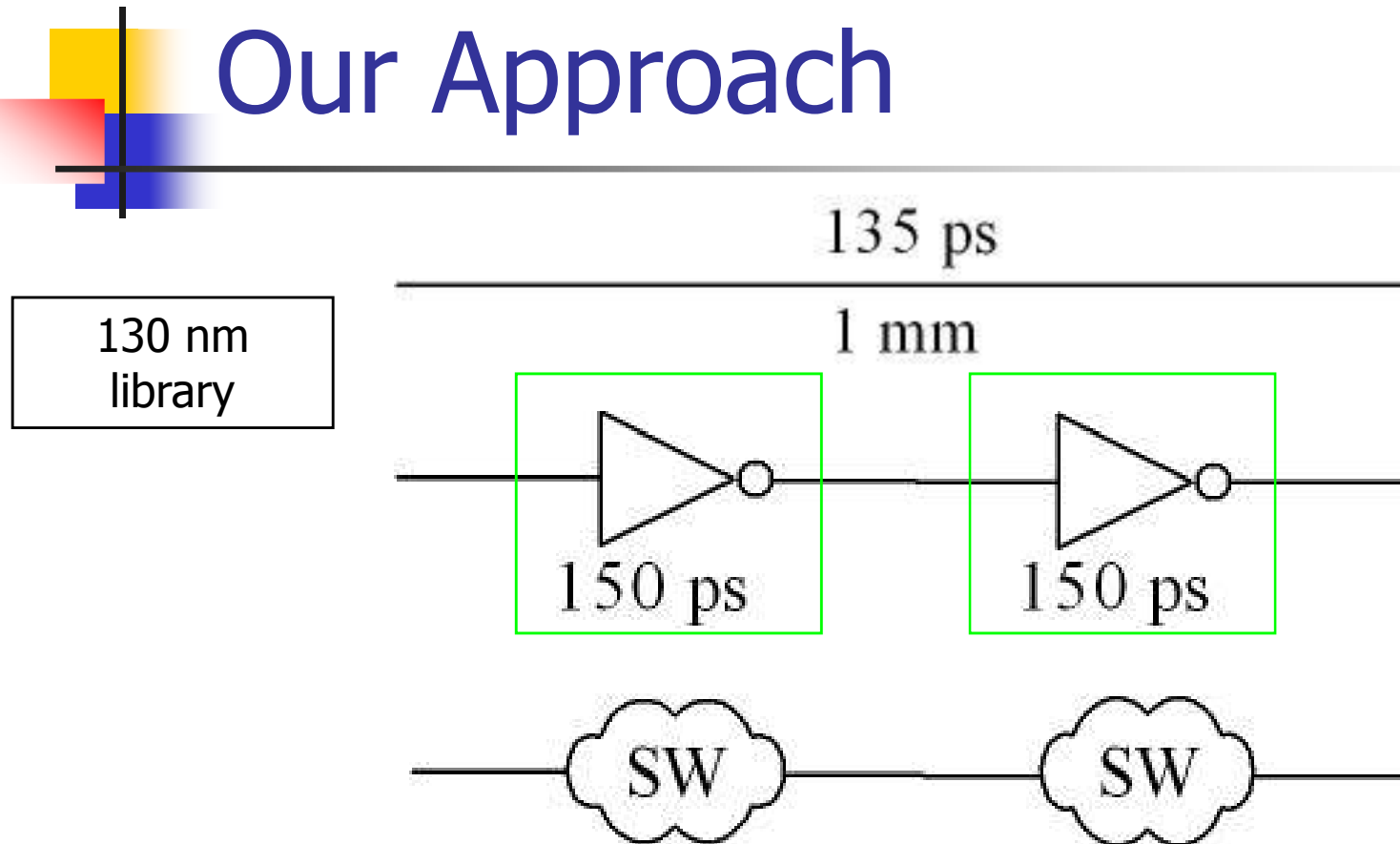
130 nm
library



Our Approach



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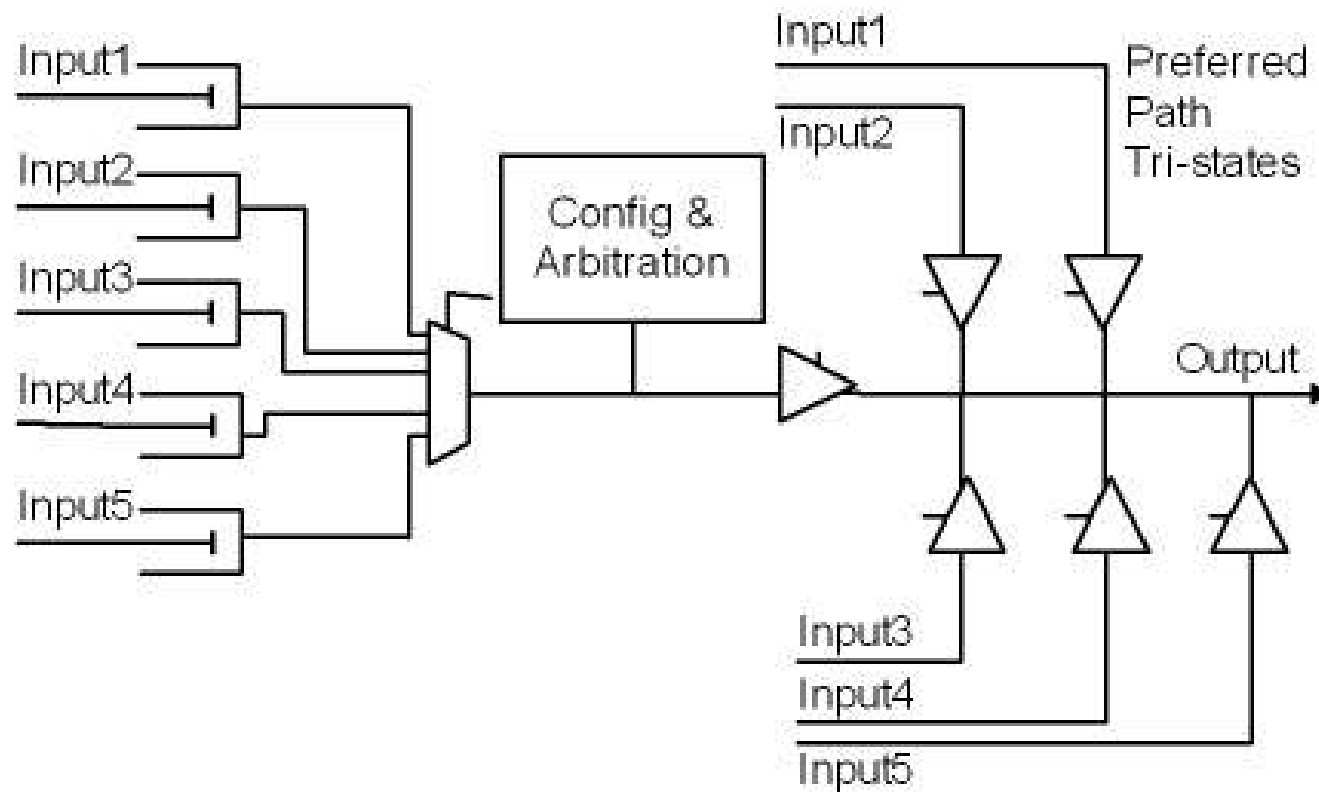
- 400 ps good scenario; 1 cycle otherwise.



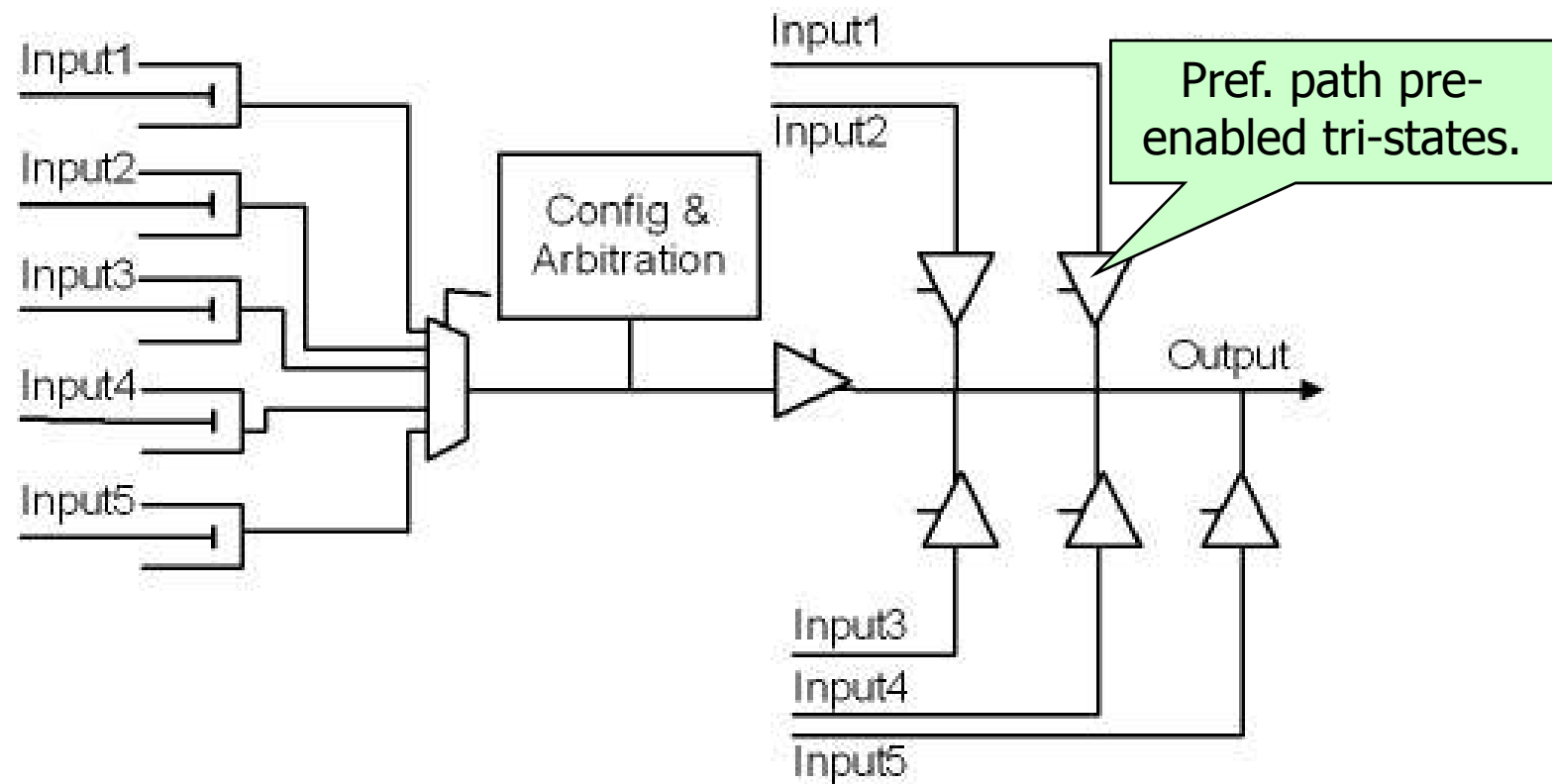
Preferred Paths

- Each output has one preferred input.
- This pref. I/O pair is connected by a single *pre-enabled* tri-state driver.
- Pre-enabling is crucial:
 - 200 ps pre-enabled mux; 500 ps otherwise.
- Later check if flits correctly forwarded.
- Thus, preferred paths are formed.
 - Reconfigurable at run-time.
 - Custom routes (shapes) allowed.

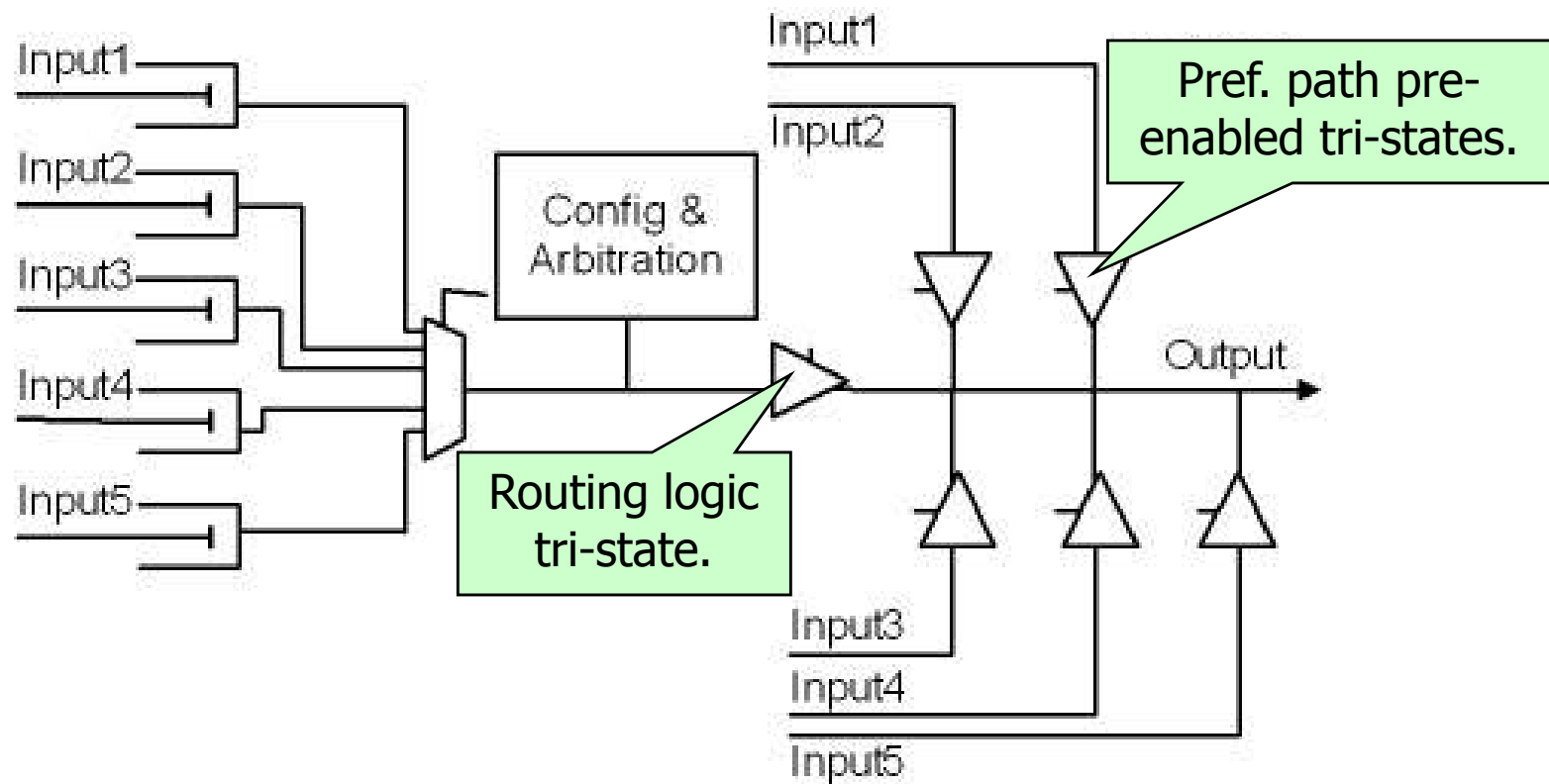
Switch Architecture - Output



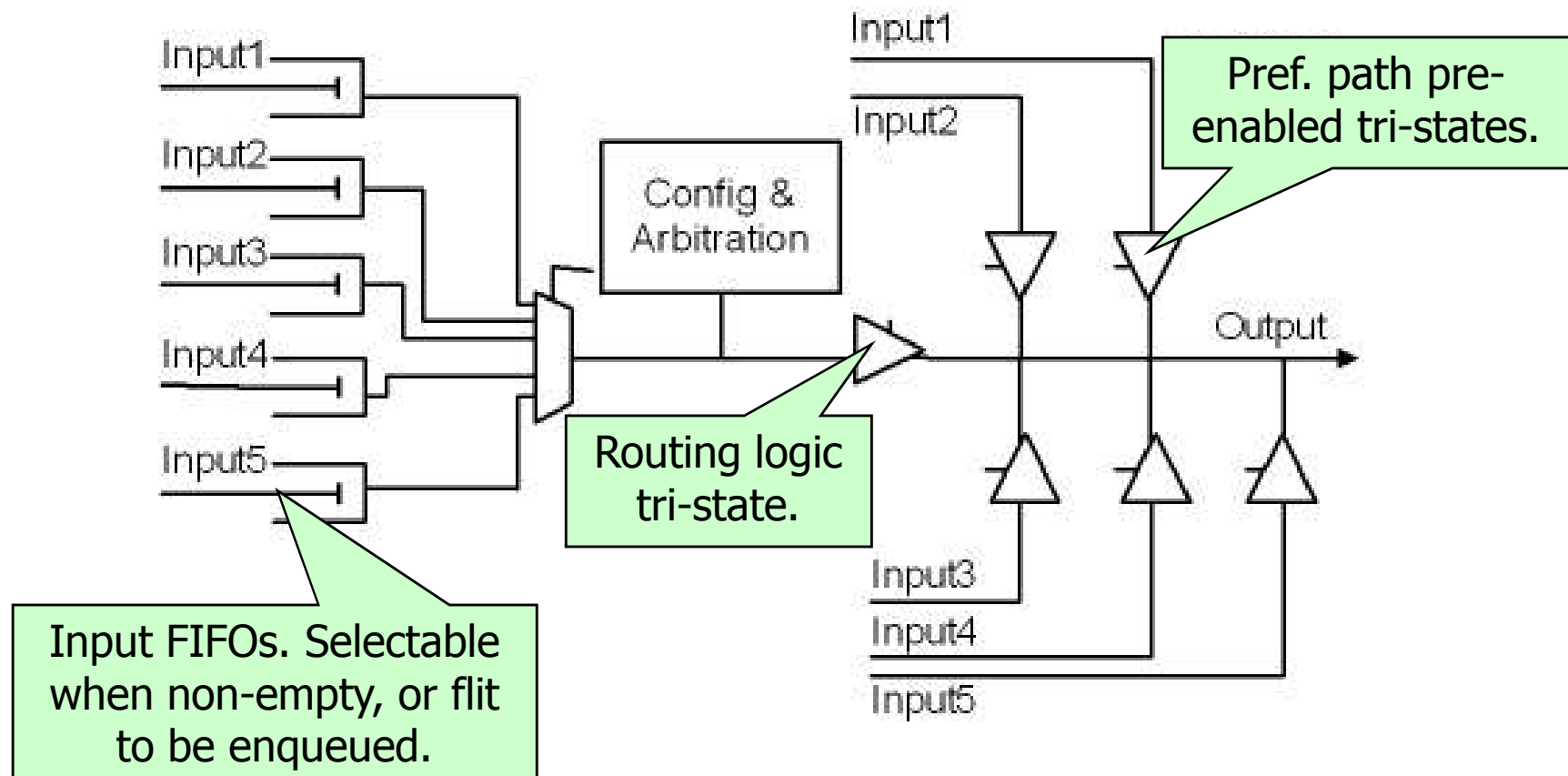
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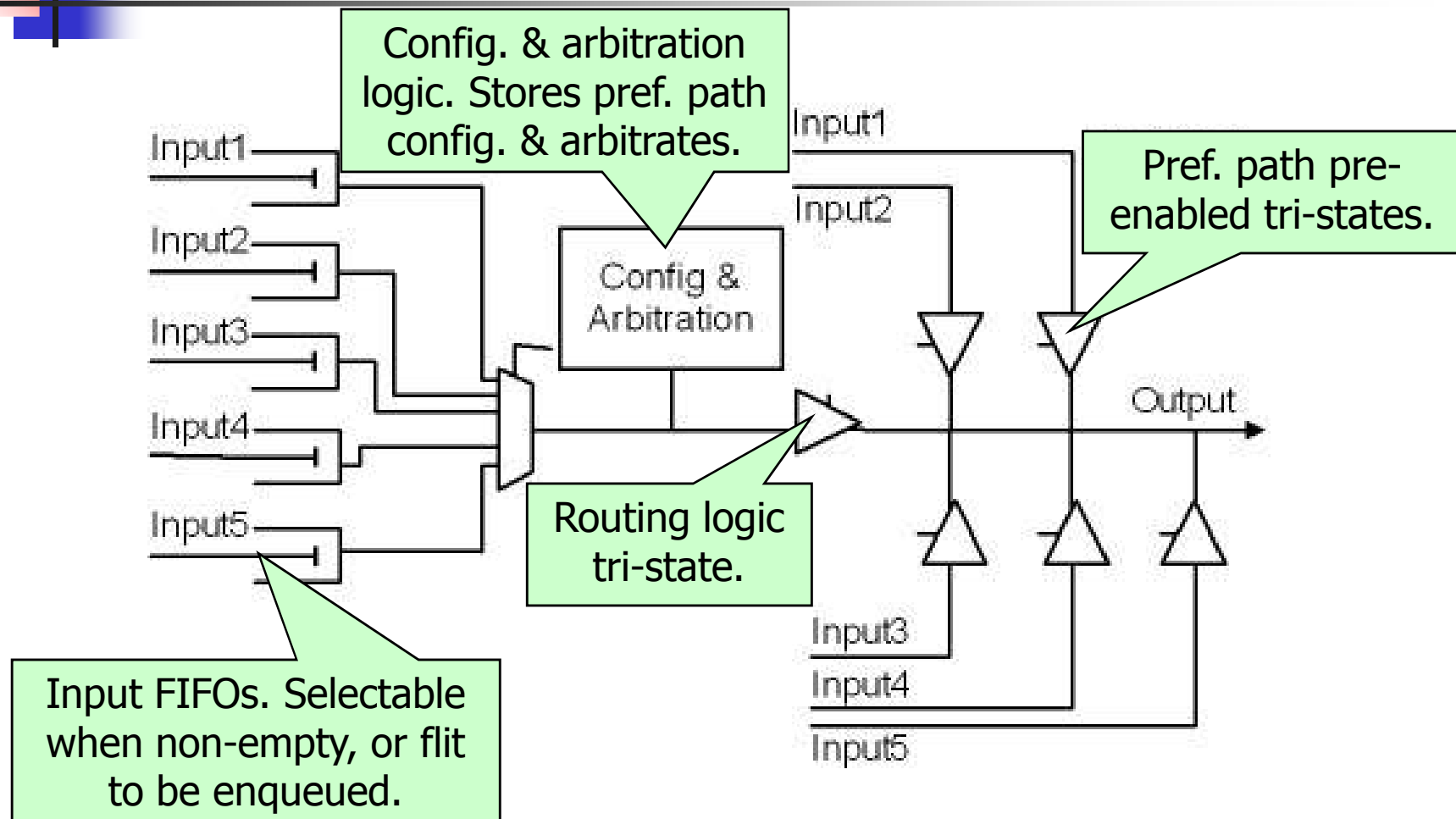
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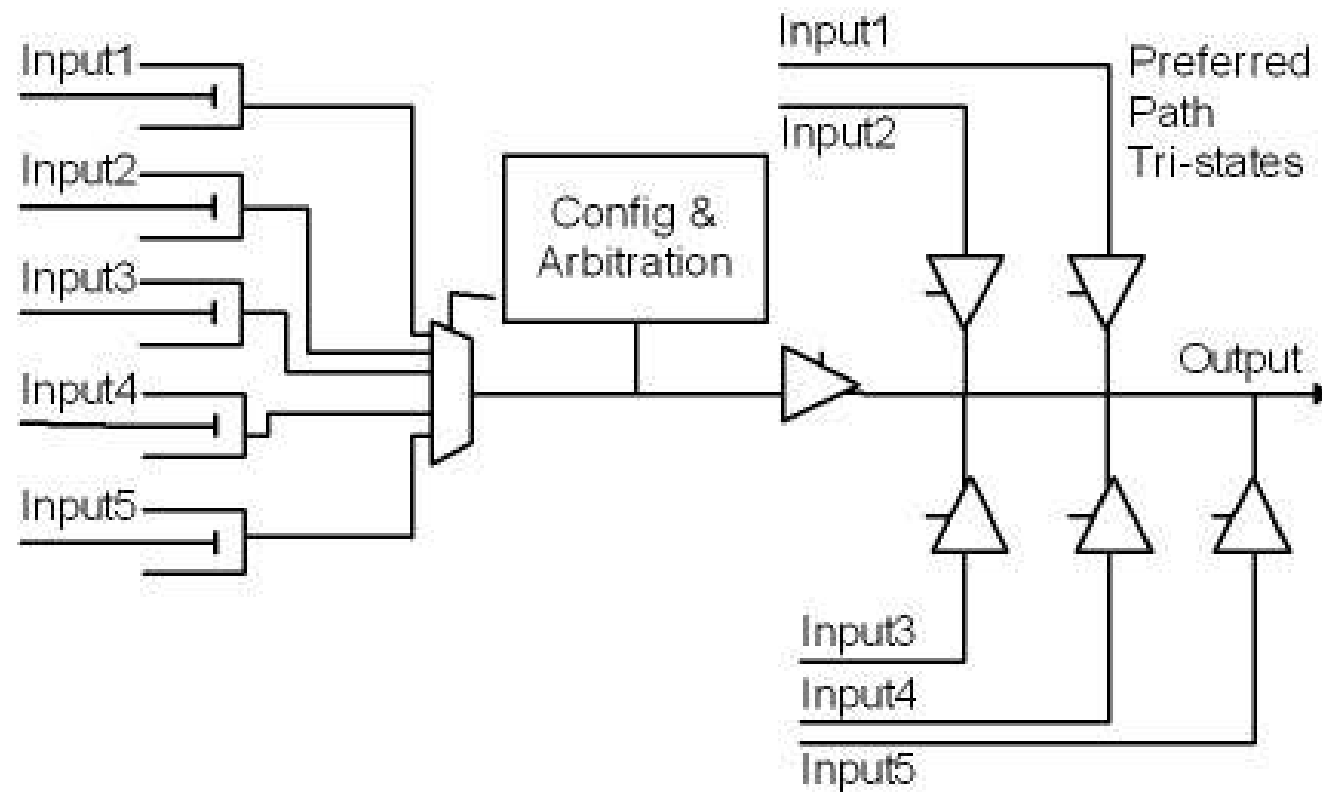
Switch Architecture - Output



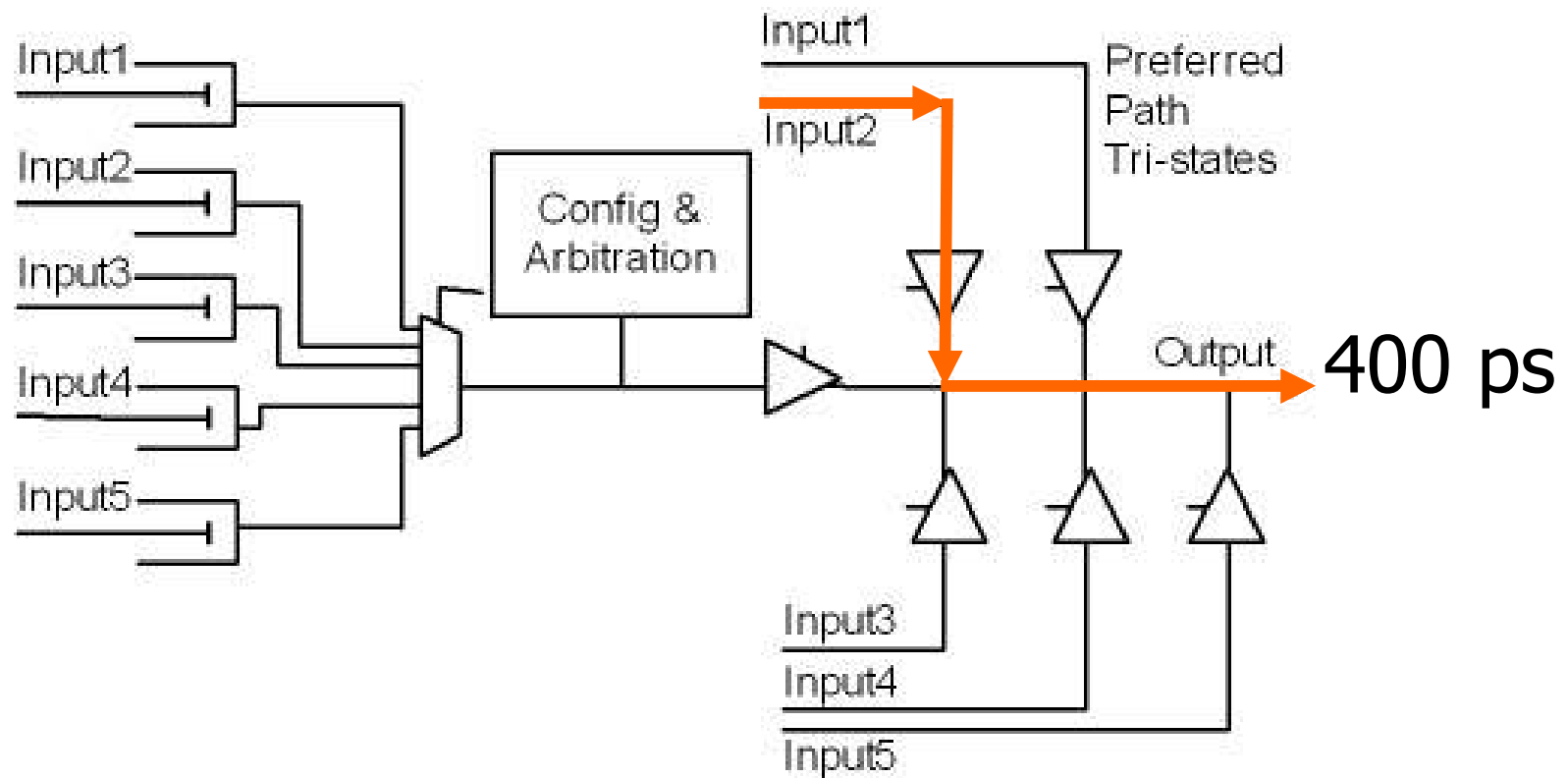
Switch Architecture - Output



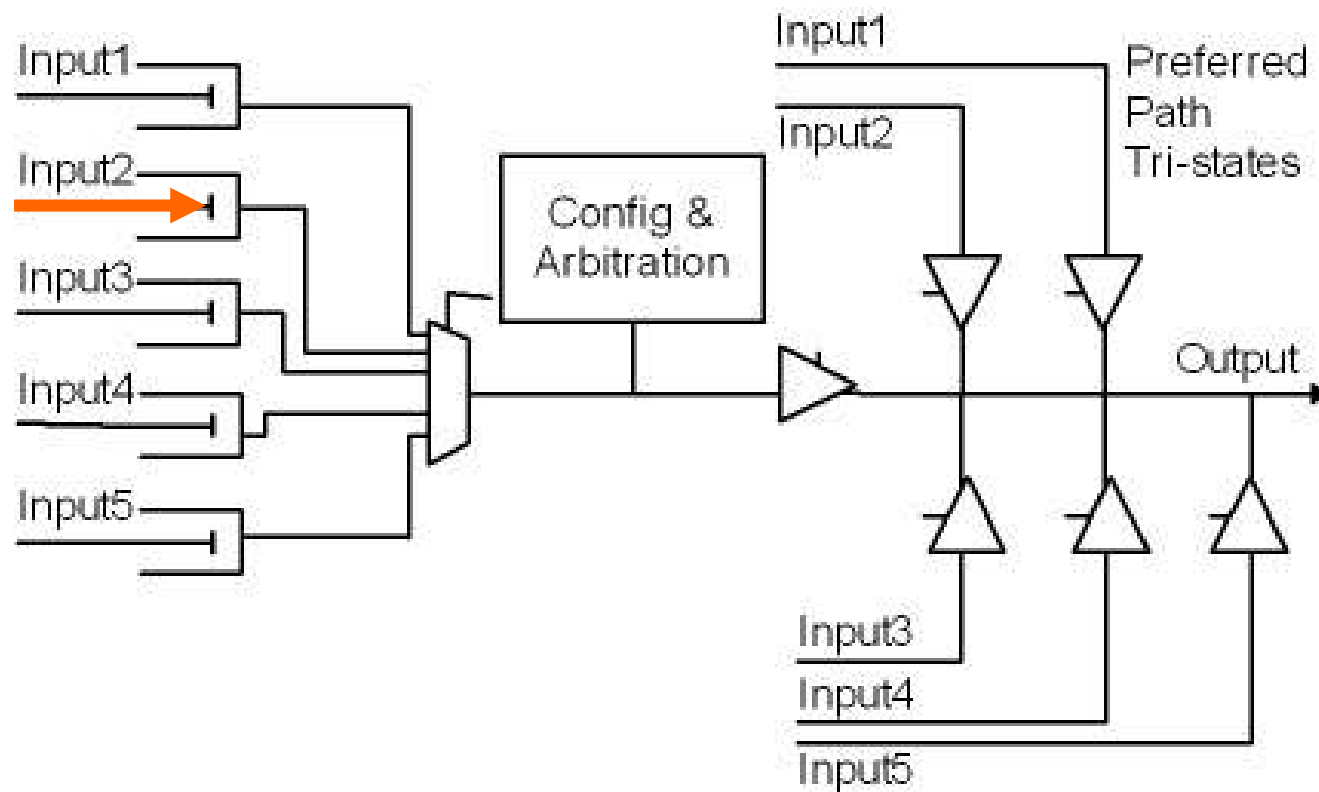
Switch Architecture - Output



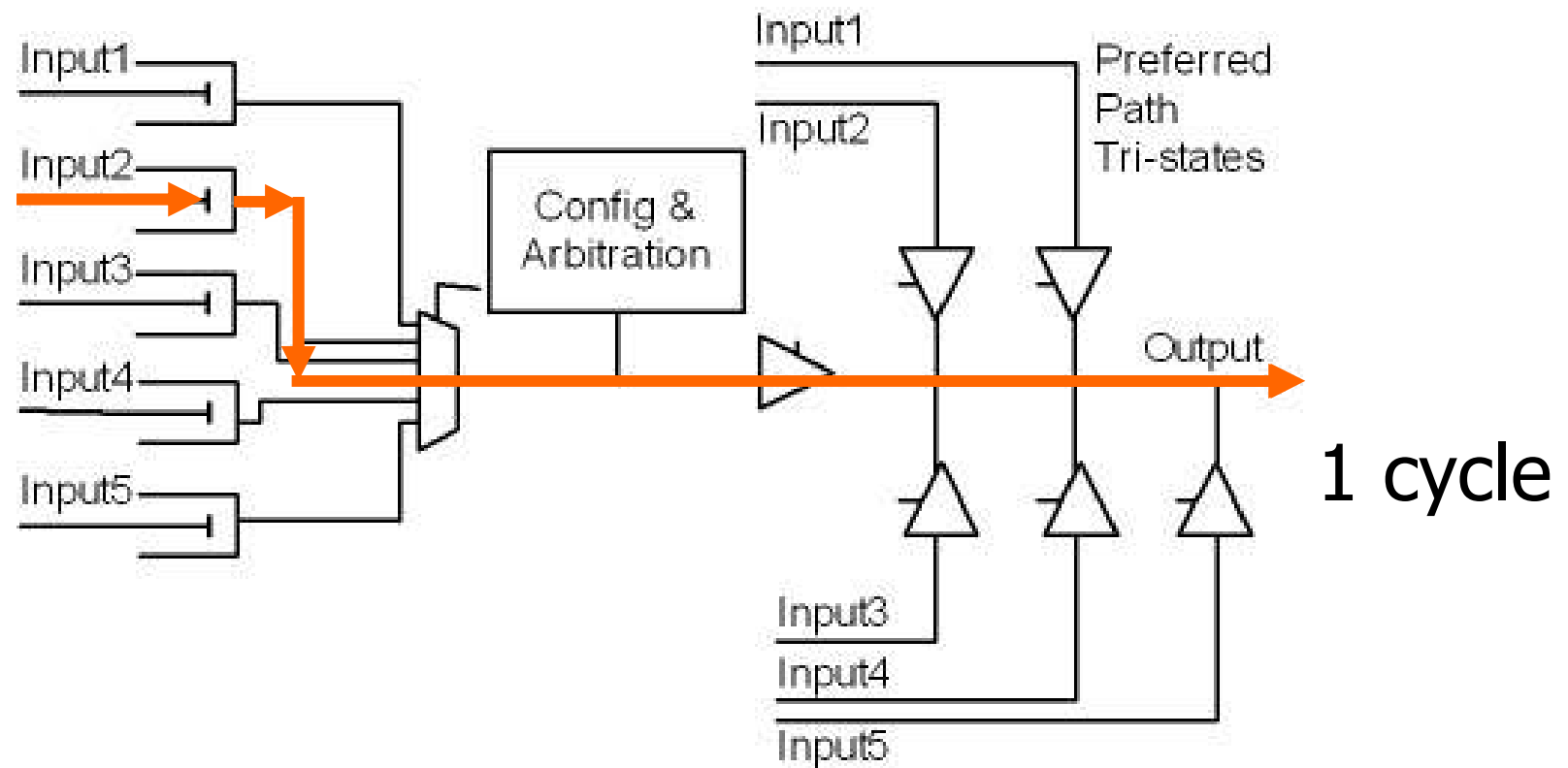
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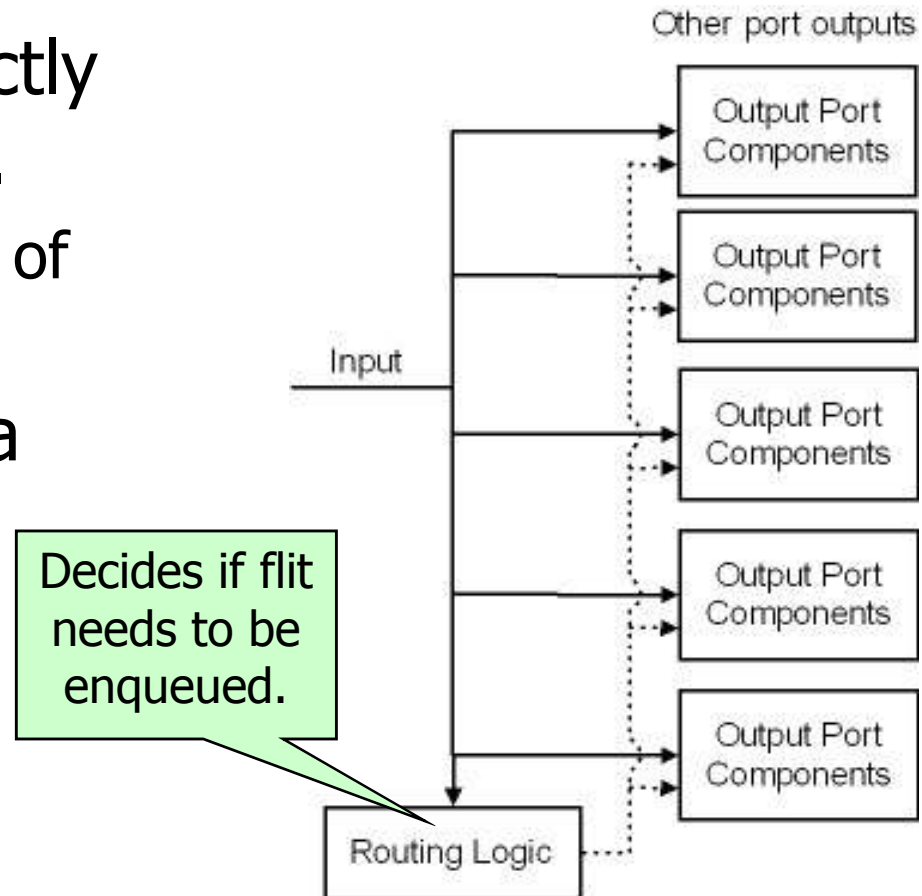


Switch Architecture - Output



Switch Architecture - Input

- Dead flits: Incorrectly eagerly forwarded.
 - Terminated at end of preferred path.
- Switch resembles a buffered crossbar.



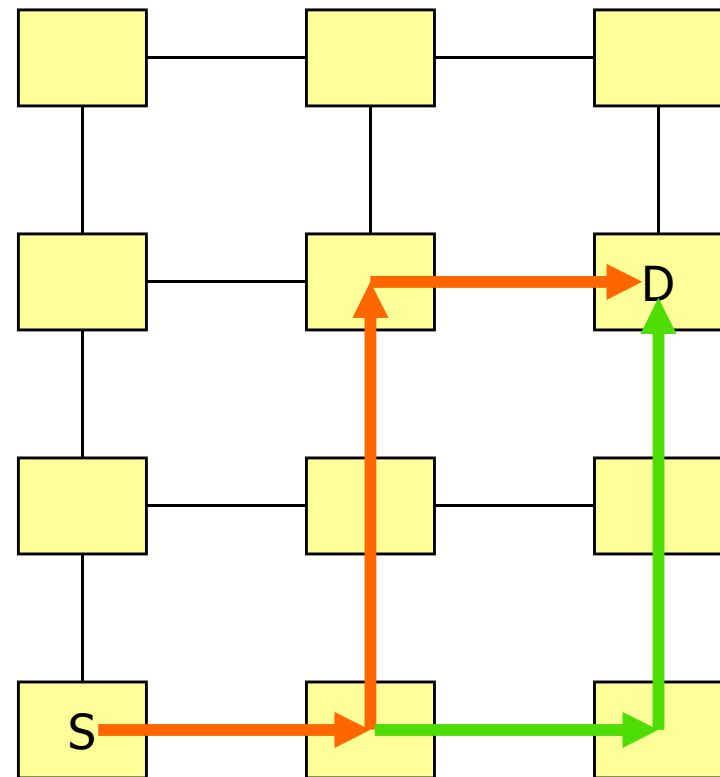


Routing Algorithm

- Deterministic routing employed.
- Non-preferred paths follow XY routing.
- We slightly modify XY routing to handle preferred paths:
 - Flit correctly eagerly forwarded if it approaches the destination in any axis.
 - Flit considered dead otherwise.

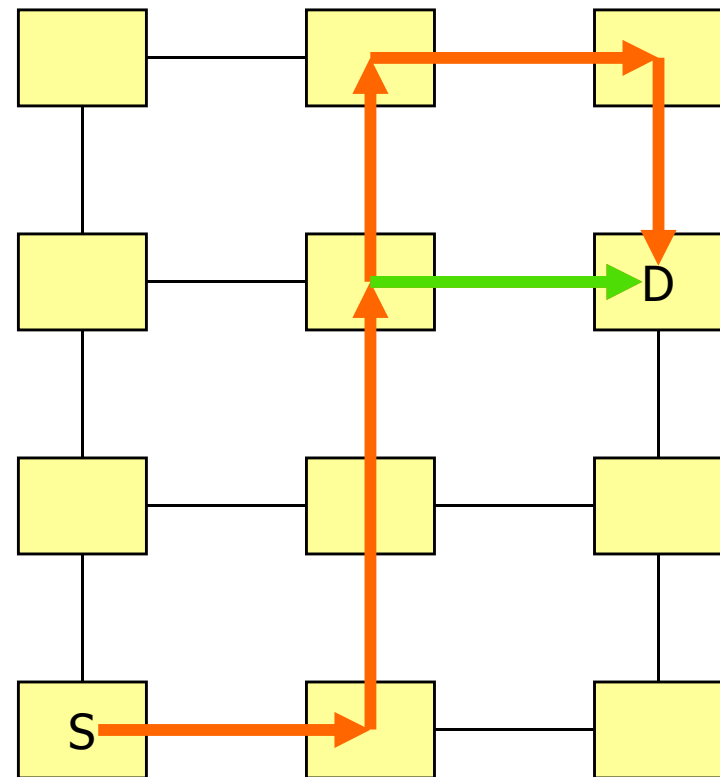
Routing Characteristics

- **Flits in preferred paths may not follow XY routing.**
- Duplicate copies of a flit may be delivered.
- XY routing.
- Pref. paths.



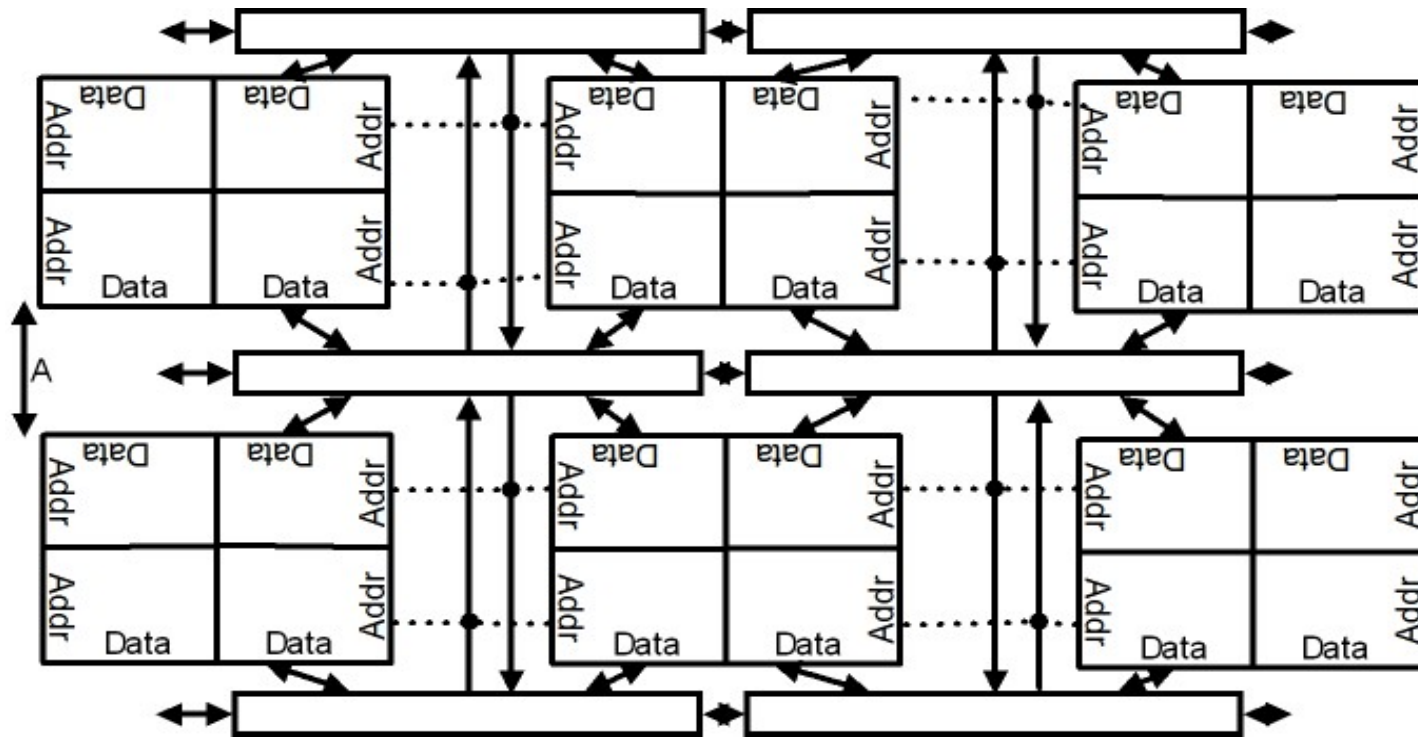
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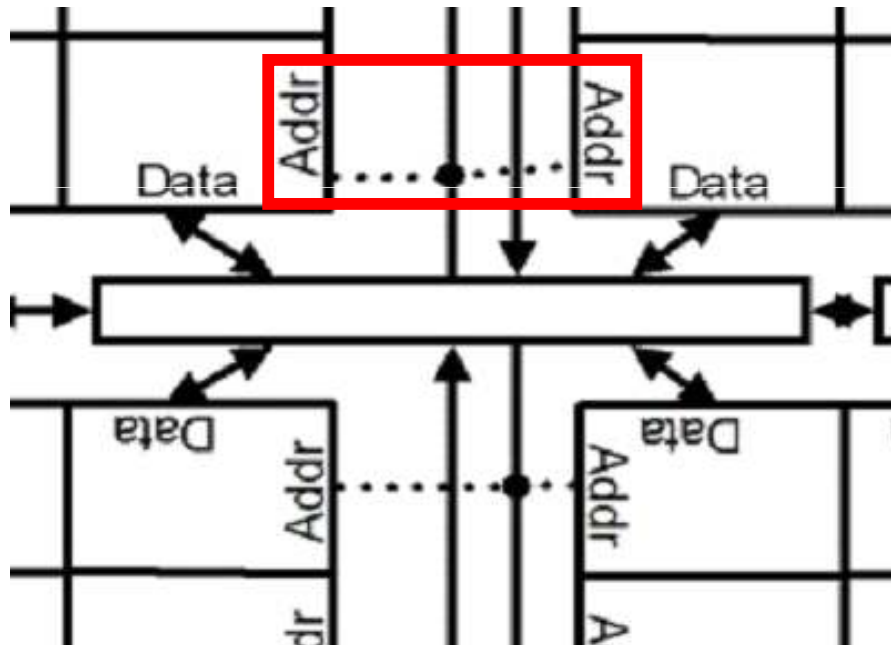


NoC Topology – Bar Floorplan

- Application: Tiled CPU and RAM blocks.
- Each switch is 6x6 and serves 4 PEs.

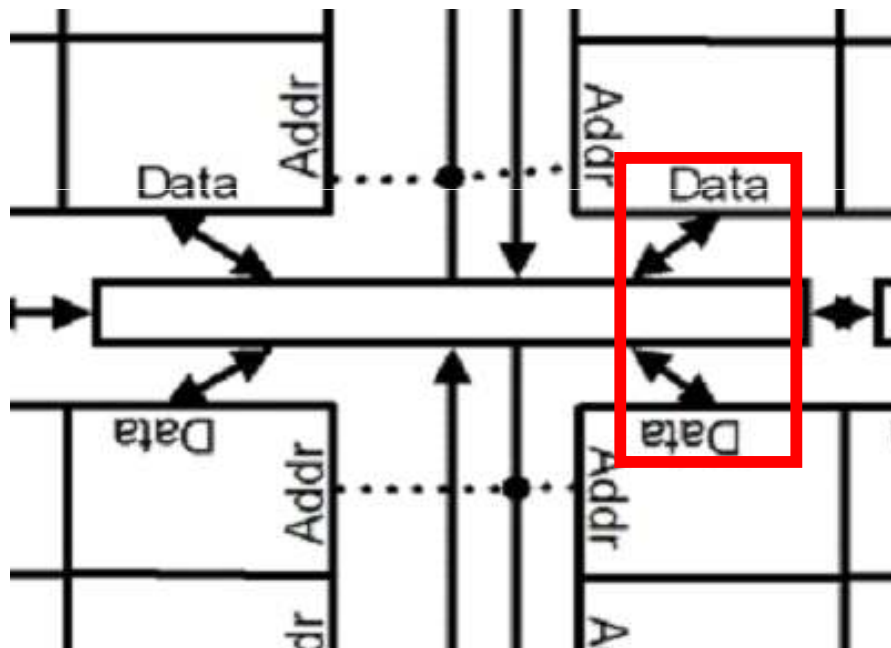


Bar Floorplan

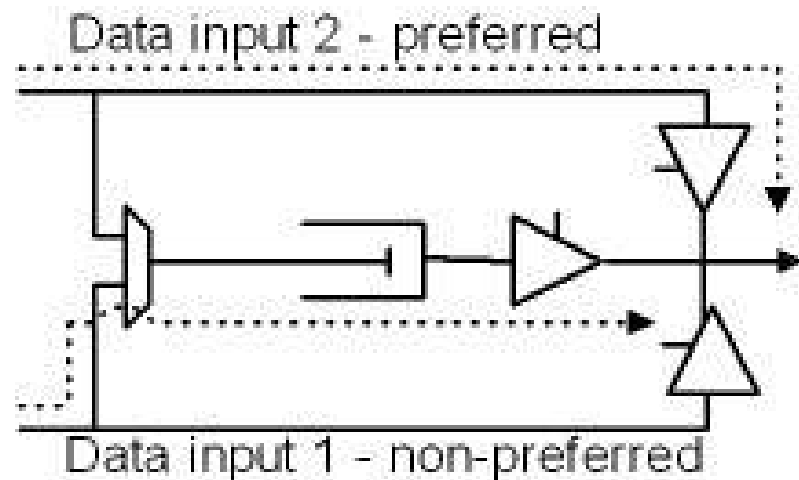


- Would be 8x12:
 - Vertical links drive address inputs.
 - 2 PE data ports served by 1 switch port.

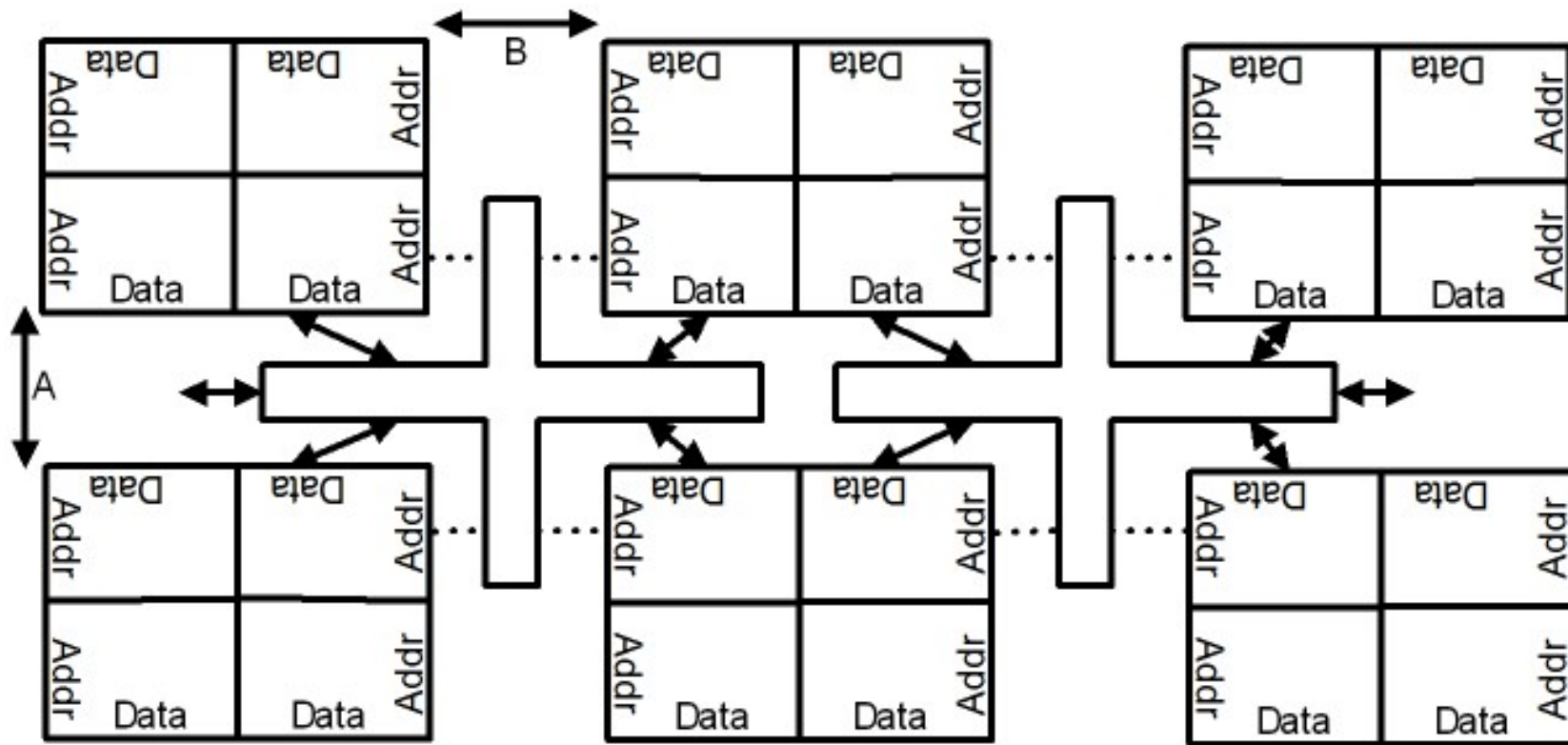
Bar Floorplan



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Cross Floorplan





Layout Results

- 130 nm implementation library. Typical case.
- Pref. path latency:
 - 300-420 ps.
 - 450-500 ps (incl. 1mm).
- 1 cycle/node otherwise.
- Past work: 1 cycle/node at 500 MHz.

| | |
|---------------------|---------|
| Clock frequency | 667 MHz |
| Flit width | 39 |
| FIFO lines | 2 |
| Number of FIFOs | 30 |
| Bar area overhead | 13% |
| Cross area overhead | 18% |
| Number of cells | 15 K |
| Number of gates | 45 K |
| Total dynamic power | 80 mW |



Advanced Issues

- Deadlock & livelock freedom.
 - Constraints to prevent circle.
 - Keep NoC functional in any case.
- Out-of-order delivery of flits in the same packet.
 - Apply reconfiguration at a “safe” time.
- Adaptive routing.



Future Work

- Synchronization issues – A flit may arrive at any time.
 - Impose preferred path constraints.
 - Implement switch asynchronously.
- Evaluation in complete system.
- Implement fault-tolerance.



Conclusion

- We approach ideal latency.
 - By pre-enabled tri-state paths.
- Our NoC is a generalized “mad-postman” [C. R. Jesshope et al, 1989].
- Our NoC is easily generalized – topology may need to be changed.
- Past NoC research can be applied for further optimizations.